

(i). All the declarative sentences to which it is possible to assign one and only one of the two possible truth values are called statements.

(ii).

$$\neg(P \vee Q) \equiv \neg P \wedge \neg Q$$

$$\neg(P \wedge Q) \equiv \neg P \vee \neg Q$$

(iii). A string of formulas is defined as

(a) Any formula is a string of formulas.

(b) If  $\alpha$  and  $\beta$  are string of formulas, then  $\alpha, \beta$  and  $\beta, \alpha$  are strings of formulas.

(c) Only those strings which are obtained by steps (a) and (b) are strings of formulas, with the exception of the empty string which is also a string of formulas.

(iv). Let  $(M, *, e_M)$  and  $(T, \Delta, e_T)$  be any two monoids. A mapping  $g: M \rightarrow T$  such that for any two elements  $a, b \in M$

$$g(a * b) = g(a) \Delta g(b)$$

and  $g(e_M) = e_T$

is called a monoid homomorphism.

(v). Let  $V$  denote a nonempty set of symbols. The set of strings over  $V$  is denoted by  $V^*$  and the set of nonempty strings by  $V^+ = V^* - \{\lambda\}$ . Then  $V^+$  is a semigroup with respect to concatenation operation. This semigroup  $(V^+, \circ)$  is called a free semigroup.

(vi). Let  $(L, \leq)$  be a lattice. For any  $a, b, c \in L$

$$a \leq c \quad \text{iff} \quad a \oplus (b * c) \leq (a \oplus b) * c$$

(vii). A lattice  $(L, *, \oplus, 0, 1)$  is said to be a complemented lattice if every element of  $L$  has at least one complement.

(viii). Let  $(B, *, \oplus, ', 0, 1)$  and  $(P, \wedge, \vee, \bar{\phantom{x}}, \alpha, \beta)$  be two Boolean algebras. A mapping  $f: B \rightarrow P$  is called a Boolean homomorphism if all the operations of the Boolean algebra are preserved, i.e. for any  $a, b \in B$

$$f(a * b) = f(a) \wedge f(b), \quad f(a \oplus b) = f(a) \vee f(b)$$
$$f(a') = \overline{f(a)}, \quad f(0) = \alpha, \quad f(1) = \beta.$$

(ix). Two Boolean forms  $\alpha(x_1, x_2, \dots, x_n)$  and  $\beta(x_1, x_2, \dots, x_n)$  are called equivalent if one can be obtained from the other by a finite number of applications of the identities of a Boolean algebra.

(\*). A context-sensitive grammar contains only productions of the form  $\alpha \rightarrow \beta$  where  $|\alpha| \leq |\beta|$ .

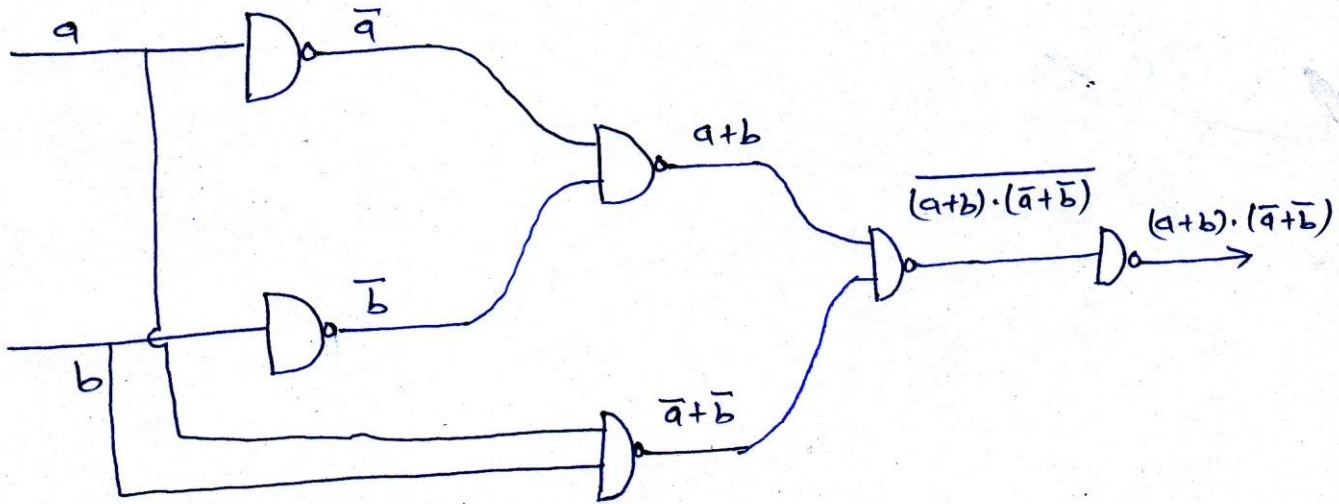
2(a). By using the equivalence of formulas, the students have to show that

$$(P \rightarrow (Q \rightarrow R)) \rightarrow ((P \rightarrow Q) \rightarrow (P \rightarrow R)) \text{ is equivalent to } T.$$

(b). By using the equivalence of formulas, the principal disjunctive normal form is obtained as

$$(\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (P \wedge Q).$$

3(a).



4(b). The students have to show that the premises  
 $P \rightarrow Q, P \rightarrow R, Q \rightarrow \neg R, P$   
will imply  $F$  (contradiction).

4(a). Let  $g$  be a semigroup homomorphism from  $(S, *)$  to  $(T, \Delta)$ . Let  $a \in S$  is an idempotent element, then

$$\begin{aligned} & a * a = a \\ & g(a * a) = g(a) \\ & g(a) \Delta g(a) = g(a) \end{aligned}$$

Hence  $g(a)$  is idempotent element.

Also let  $a, b \in S$  commute with each other

$$\text{ie. } a * b = b * a$$

$$\text{then } g(a * b) = g(b * a)$$

$$g(a) \Delta g(b) = g(b) \Delta g(a).$$

then  $g(a)$  and  $g(b)$  commute with each other.

(b). Let  $(S, *)$  be a finite semigroup and  $a \in S$ .

then  $a^n \in S \quad \forall n \in \mathbb{N}$

but  $S$  is finite, so  $\exists r \neq s \in \mathbb{N}$  s.t.  $s > r$   
(ie  $s = r + k$ )

and  $a^r = a^s$

$$a^r = a^{r+k}$$

then  $a^{2r+k} = a^r \cdot a^{r+k} = a^r \cdot a^r = a^{2r}$

Similarly  $a^{mr+k} = a^{mr} \quad \forall m \in \mathbb{N}$ .

Again  $a^{mr+2k} = a^{mr+k} \cdot a^k = a^{mr} \cdot a^k = a^{mr+k} = a^{mr}$

Similarly  $a^{mr+nk} = a^{mr} \quad \forall n \in \mathbb{N}$ .

Put  $m = k$  and  $n = r$ , we get

$$a^{2kr} = a^{kr}$$

ie  $(a^{kr})^2 = a^{kr}$

Put  $y = a^{kr}$  then

$$y^2 = y$$

ie  $y = a^{kr} \in S$  is an idempotent element.

5(d). Let  $(L, \leq)$  be a lattice and  $a, b, c \in L$ .

The students have to prove

$$b \leq c \implies \begin{cases} a * b \leq a * c \\ a \oplus b \leq a \oplus c \end{cases}$$

(b).

Since  $a * b \leq a \leq a \oplus c$

and  $a * b \leq b \leq b \oplus d$

then  $a * b \leq (a \oplus c) * (b \oplus d) \quad \text{--- ①}$

Also  $c * d \leq c \leq a \oplus c$

and  $c * d \leq d \leq b \oplus d$

then  $c * d \leq (a \oplus c) * (b \oplus d)$  — ②

From ① & ②

$$(a * b) \oplus (c * d) \leq (a \oplus c) * (b \oplus d).$$

### 6. Lattice as a poset.

A lattice is a partially ordered set  $(L, \leq)$  in which every pair of elements  $a, b \in L$  has greatest lower bound and a least upper bound.

### Lattice as algebraic structure

A lattice is an algebraic system  $(L, *, \oplus)$  with two binary operations  $*$  and  $\oplus$  on  $L$  which are both, commutative, associative and satisfy the absorption laws.

The students have to obtain the properties of definition 2 from def. 1 and also from def. 1 to def. 2.

7.(a). The sum of product canonical form is obtained as

$$\begin{aligned} & x_1' x_2' x_3 x_4 + x_1' x_2' x_3 x_4' + x_1' x_2' x_3' x_4 + x_1' x_2' x_3' x_4' \\ & + x_1 x_2 x_3 x_4 + x_1' x_2 x_3 x_4 + x_1 x_2' x_3 x_4 \end{aligned}$$

(b). Simplified expression is as

$$(a * b') \oplus (b' * c)$$

8. Students have to discuss Context-sensitive grammar, Context-free grammar, regular grammar and free grammar.